

# Codes and invariant theory.

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Clifford Weil Groups and Symmetrizations



# Complete weight enumerators,

Let  $V$  be a finite abelian group (e.g.  $V = \mathbb{F}_q$ ) and  $C \subseteq V^N$ . For  $c = (c_1, \dots, c_N) \in V^N$  and  $v \in V$  put

$$a_v(c) := |\{i \in \{1, \dots, N\} \mid c_i = v\}|.$$

Then

$$\text{cwe}_C := \sum_{c \in C} \prod_{v \in V} x_v^{a_v(c)} \in \mathbb{C}[x_v : v \in V]$$

is called the **complete weight enumerator** of  $C$ .

## The tetracode.

$$t_4 := \begin{bmatrix} 1 & 1 & 1 & 0 \\ 0 & 1 & 2 & 1 \end{bmatrix} \leq \mathbb{F}_3^4$$

$$\text{cwe}_{t_4}(x_0, x_1, x_2) = x_0^4 + x_0 x_1^3 + x_0 x_2^3 + 3x_0 x_1^2 x_2 + 3x_0 x_1 x_2^2.$$

$$\text{hwe}_{t_4}(x, y) = \text{cwe}_{t_4}(x, y, y) = x^4 + 8xy^3.$$

**Clear.**  $\text{hwe}_C(x, y) = \text{cwe}_C(x, y, \dots, y)$

# A formal notion of a Type of a code.

## Definition of Type, part I

A **Type** is a quadrupel  $(R, V, \Phi, \beta)$  with

- ▶  $R$  is a finite ring (with 1) and  $J : R \rightarrow R$  an involution of  $R$ .  
 $(ab)^J = b^J a^J$  and  $(a^J)^J = a$  for all  $a, b \in R$
- ▶  $V$  a finite left  $R$ -module.
- ▶  $\beta : V \times V \rightarrow \mathbb{Q}/\mathbb{Z}$  regular,  $\epsilon$ -hermitian:  
 $\beta(rv, w) = \beta(v, r^J w)$  for  $r \in R, v, w \in V$ ,  
 $v \mapsto \beta(v, \cdot) \in \text{Hom}(V, \mathbb{Q}/\mathbb{Z})$  isomorphism,  
 $\epsilon \in Z(R)$ ,  $\epsilon \epsilon^J = 1$   $\beta(v, w) = \beta(w, \epsilon v)$  for  $v, w \in V$ .
- ▶  $\Phi \subset \text{Quad}_0(V, \mathbb{Q}/\mathbb{Z})$  a set of quadratic mappings on  $V$ .

and certain additional properties.

# Codes of a given Type.

Let  $(R, V, \Phi, \beta)$  be a Type.

## Definition.

- ▶ A **code**  $C$  over the alphabet  $V$  is an  $R$ -submodule of  $V^N$ .
- ▶ The **dual code** (with respect to  $\beta$ ) is

$$C^\perp := \{x \in V^N \mid \beta^N(x, c) = \sum_{i=1}^N \beta(x_i, c_i) = 0 \text{ for all } c \in C\}.$$

$C$  is called **self-dual** (with respect to  $\beta$ ) if  $C = C^\perp$ .

- ▶ Then  $C$  is called **isotropic** (with respect to  $\Phi$ ) if

$$\phi^N(c) := \sum_{i=1}^N \phi(c_i) = 0 \text{ for all } c \in C \text{ and } \phi \in \Phi.$$

# A formal notion of a Type of a code.

## Definition

The quadruple  $(R, V, \Phi, \beta)$  as above is called a **Type** if

- $\Phi \leq \text{Quad}_0(V, \mathbb{Q}/\mathbb{Z})$  is a subgroup and for all  $r \in R, \phi \in \Phi$  the mapping  $\phi[r] : x \mapsto \phi(rx)$  is again in  $\Phi$ .

Then  $\Phi$  is an *R-qmodule*.

- For all  $\phi \in \Phi$  there is some  $r_\phi \in R$  such that

$$\lambda(\phi)(v, w) = \phi(v + w) - \phi(v) - \phi(w) = \beta(v, r_\phi w) \text{ for all } v, w \in V.$$

- For all  $r \in R$  the mapping

$$\phi_r : V \rightarrow \mathbb{Q}/\mathbb{Z}, v \mapsto \beta(v, rv) \text{ lies in } \Phi.$$

# Type I,II,III,IV in the new language.

## Type I codes ( $2_I$ )

$$R = \mathbb{F}_2 = V, \beta(x, y) = \frac{1}{2}xy, \Phi = \{\varphi : x \mapsto \frac{1}{2}x^2 = \beta(x, x), 0\}$$

## Type II code ( $2_{II}$ ).

$$R = \mathbb{F}_2 = V, \beta(x, y) = \frac{1}{2}xy, \Phi = \{\phi : x \mapsto \frac{1}{4}x^2, 2\phi = \varphi, 3\phi, 0\}$$

## Type III codes (3).

$$R = \mathbb{F}_3 = V, \beta(x, y) = \frac{1}{3}xy, \Phi = \{\varphi : x \mapsto \frac{1}{3}x^2 = \beta(x, x), 2\varphi, 0\}$$

## Type IV codes ( $4^H$ ).

$$R = \mathbb{F}_4 = V, \beta(x, y) = \frac{1}{2} \text{tr}(x\bar{y}), \Phi = \{\varphi : x \mapsto \frac{1}{2}x\bar{x}, 0\}$$

where  $\bar{x} = x^2$ .

# The Clifford-Weil group associated to a Type.

## Definition.

Let  $T := (R, V, \beta, \Phi)$  be a Type. Then the associated **Clifford-Weil group**  $\mathcal{C}(T)$  is a subgroup of  $\mathrm{GL}_{|V|}(\mathbb{C})$

$$\mathcal{C}(T) = \langle m_r, d_\phi, h_{e, u_e, v_e} \mid r \in R^*, \phi \in \Phi, e = u_e v_e \in R \text{ sym. id. } \rangle$$

Let  $(e_v \mid v \in V)$  denote a basis of  $\mathbb{C}^{|V|}$ . Then

$$m_r : e_v \mapsto e_{rv}, \quad d_\phi : e_v \mapsto \exp(2\pi i \phi(v)) e_v$$

$$h_{e, u_e, v_e} : e_v \mapsto |eV|^{-1/2} \sum_{w \in eV} \exp(2\pi i \beta(w, v_e v)) e_{w + (1-e)v}$$

# Invariance of complete weight enumerators.

## Theorem.

Let  $C \leq V^N$  be a self-dual isotropic code of Type  $T$ . Then  $\text{cwe}_C$  is invariant under  $\mathcal{C}(T)$ .

## Proof.

Invariance under  $m_r$  ( $r \in R^*$ ) because  $C$  is a code.

Invariance under  $d_\phi$  ( $\phi \in \Phi$ ) because  $C$  is isotropic.

Invariance under  $h_{e,u_e,v_e}$  because  $C$  is self dual.

## The main theorem.(N., Rains, Sloane (1999-2006))

If  $R$  is a direct product of matrix rings over chain rings, then

$$\text{Inv}(\mathcal{C}(T)) = \langle \text{cwe}_C \mid C \text{ of Type } T \rangle.$$

# The Clifford-Weil groups for Type I and II.

## Type I codes ( $2_I$ )

$$R = \mathbb{F}_2 = V, \beta(x, y) = \frac{1}{2}xy, \Phi = \{\varphi : x \mapsto \frac{1}{2}x^2 = \beta(x, x), 0\}$$

$$\mathcal{C}(I) = \langle d_\varphi = \text{diag}(1, -1), h_{1,1,1} = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix} = h_2 \rangle = G_I$$

## Type II codes ( $2_{II}$ ).

$$R = \mathbb{F}_2 = V, \beta(x, y) = \frac{1}{2}xy, \Phi = \{\phi : x \mapsto \frac{1}{4}x^2, 2\phi = \varphi, 3\phi, 0\}$$

$$\mathcal{C}(II) = \langle d_\phi = \text{diag}(1, i), h_2 \rangle = G_{II}$$

# The Clifford-Weil groups for Type III and IV.

Type III codes (3).

$$R = \mathbb{F}_3 = V, \beta(x, y) = \frac{1}{3}xy, \Phi = \{\varphi : x \mapsto \frac{1}{3}x^2 = \beta(x, x), 2\varphi, 0\}$$

$$\mathcal{C}(\text{III}) = \langle m_2 = \begin{pmatrix} 100 \\ 001 \\ 010 \end{pmatrix}, d_\varphi = \text{diag}(1, \zeta_3, \zeta_3), h_{1,1,1} = \frac{1}{\sqrt{3}} \begin{pmatrix} 1 & 1 & 1 \\ 1 & \zeta_3 & \zeta_3^2 \\ 1 & \zeta_3^2 & \zeta_3 \end{pmatrix} \rangle$$

Type IV codes ( $4^H$ ).

$$R = \mathbb{F}_4 = V, \beta(x, y) = \frac{1}{2} \text{tr}(x\bar{y}), \Phi = \{\varphi : x \mapsto \frac{1}{2}x\bar{x}, 0\}$$

$$\mathcal{C}(\text{IV}) = \langle m_\omega = \begin{pmatrix} 1000 \\ 0001 \\ 0100 \\ 0010 \end{pmatrix}, d_\varphi = \text{diag}(1, -1, -1, -1), h_{1,1,1} = \frac{1}{2} \begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & 1 & -1 & -1 \\ 1 & -1 & 1 & -1 \\ 1 & -1 & -1 & 1 \end{pmatrix} \rangle$$



# Symmetrizations.

## Definition

Let  $(R, J)$  be a ring with involution. Then the **central unitary group** is

$$\mathrm{ZU}(R, J) := \{g \in Z(R) \mid gg^J = g^Jg = 1\}.$$

## Theorem.

Let  $T = (R, V, \beta, \Phi)$  be a Type and

$$U := \{u \in \mathrm{ZU}(R, J) \mid \phi(uv) = \phi(v) \text{ for all } \phi \in \Phi, v \in V\}.$$

Then  $m(U) := \{m_u \mid u \in U\}$  is in the center of  $\mathcal{C}(T)$ .

## Example.

$R = \mathbb{F}_2$  or  $R = \mathbb{F}_3$  then  $\mathrm{ZU}(R, \mathrm{id}) = R - \{0\}$ .

If  $R = \mathbb{F}_4$  then  $\mathrm{ZU}(R, \mathrm{id}) = \{1\}$ , but  $\mathrm{ZU}(R, \neg) = R - \{0\}$ .

# Symmetrized Clifford-Weil groups.

## Definition.

Let  $U \leq \mathrm{ZU}(R, J)$  and  $X_0, \dots, X_n$  be the  $U$ -orbits on  $V$ .

The  $U$ -symmetrized Clifford-Weil group is

$$\mathcal{C}^{(U)}(T) = \{g^{(U)} \mid g \in \mathcal{C}(T)\} \leq \mathrm{GL}_{n+1}(\mathbb{C})$$

If

$$g\left(\frac{1}{|X_i|} \sum_{v \in X_i} e_v\right) = \sum_{j=0}^n a_{ij} \left(\frac{1}{|X_j|} \sum_{w \in X_j} e_w\right)$$

then

$$g^{(U)}(x_i) = \sum_{j=0}^n a_{ij} x_j.$$

## Remark.

The invariant ring of  $\mathcal{C}^{(U)}(T)$  consists of the  $U$ -symmetrized invariants of  $\mathcal{C}(T)$ .

# Symmetrized weight enumerators.

## Definition.

Let  $U$  permute the elements of  $V$  and let  $C \leq V^N$ . Let  $X_0, \dots, X_n$  denote the orbits on  $U$  on  $V$  and for  $c = (c_1, \dots, c_N) \in C$  and  $0 \leq j \leq n$  define

$$a_j(c) = |\{1 \leq i \leq N \mid c_i \in X_j\}|$$

Then the  $U$ -symmetrized weight-enumerator of  $C$  is

$$\text{cwe}_C^{(U)} = \sum_{c \in C} \prod_{j=0}^n x_j^{a_j(c)} \in \mathbb{C}[x_0, \dots, x_n]$$

## Remark.

If the invariant ring of  $\mathcal{C}(T)$  is spanned by the complete weight enumerators of self-dual codes of Type  $T$ , then the invariant ring of  $\mathcal{C}^{(U)}(T)$  is spanned by the  $U$ -symmetrized weight-enumerators of self-dual codes of Type  $T$ .

# Gleason's Theorem revisited.

## Remark

For Type I,II,III,IV the central unitary group  $ZU(R, J)$  is transitive on  $V - \{0\}$ , so there are only two orbits:

$$x \leftrightarrow \{0\}, \quad y \leftrightarrow V - \{0\}$$

and the symmetrized weight enumerators are the Hamming weight enumerators.

$$\mathcal{C}(\text{III}) = \langle m_2 = \begin{pmatrix} 100 \\ 001 \\ 010 \end{pmatrix}, d_\varphi = \text{diag}(1, \zeta_3, \zeta_3), h_{1,1,1} = \frac{1}{\sqrt{3}} \begin{pmatrix} 1 & 1 & 1 \\ 1 & \zeta_3 & \zeta_3^2 \\ 1 & \zeta_3^2 & \zeta_3 \end{pmatrix} \rangle$$

yields the symmetrized Clifford-Weil group  $G_{\text{III}} = \mathcal{C}^{(U)}(\text{III})$

$$\mathcal{C}^{(U)}(\text{III}) = \langle m_2^{(U)} = I_2, d_\varphi^{(U)} = \text{diag}(1, \zeta_3), h_{1,1,1}^{(U)} = h_3 = \frac{1}{\sqrt{3}} \begin{pmatrix} 1 & 2 \\ 1 & -1 \end{pmatrix} \rangle$$

# The symmetrized Clifford-Weil group of Type IV.

$$\mathcal{C}(\text{IV}) = \langle m_\omega = \begin{pmatrix} 1000 \\ 0001 \\ 0100 \\ 0010 \end{pmatrix}, d_\varphi = \text{diag}(1, -1, -1, -1), h_{1,1,1} = \frac{1}{2} \begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & 1 & -1 & -1 \\ 1 & -1 & 1 & -1 \\ 1 & -1 & -1 & 1 \end{pmatrix} \rangle$$

yields the symmetrized Clifford-Weil group  $G_{\text{IV}} = \mathcal{C}^{(U)}(\text{IV})$

$$\mathcal{C}^{(U)}(\text{IV}) = \langle m_\omega^{(U)} = I_2, d_\varphi^{(U)} = \text{diag}(1, -1), h_{1,1,1}^{(U)} = h_4 = \frac{1}{2} \begin{pmatrix} 1 & 3 \\ 1 & -1 \end{pmatrix} \rangle$$

# Hermitian codes over $\mathbb{F}_9$

$$(9^H) : R = V = \mathbb{F}_9, \beta(x, y) = \frac{1}{3} \operatorname{tr}(x\bar{y}), \Phi = \{\varphi : x \mapsto \frac{1}{3}x\bar{x}, 2\varphi, 0\}.$$

Let  $\alpha$  be a primitive element of  $\mathbb{F}_9$  and put  $\zeta = \zeta_3 \in \mathbb{C}$ . Then with respect to the  $\mathbb{C}$ -basis

$$(0, 1, \alpha, \alpha^2, \alpha^3, \alpha^4, \alpha^5, \alpha^6, \alpha^7)$$

of  $\mathbb{C}[V]$ , the associated Clifford-Weil group  $\mathcal{C}(9^H)$  is generated by  $d_\varphi := \operatorname{diag}(1, \zeta, \zeta^2, \zeta, \zeta^2, \zeta, \zeta^2, \zeta, \zeta^2)$ ,

$$m_\alpha := \begin{pmatrix} 100000000 \\ 000000001 \\ 010000000 \\ 001000000 \\ 000100000 \\ 000010000 \\ 000001000 \\ 000000100 \\ 000000010 \end{pmatrix}, \quad h := \frac{1}{3} \begin{pmatrix} 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 1\zeta^2 & \zeta & 1 & \zeta & \zeta & \zeta^2 & 1 & \zeta^2 \\ 1 & \zeta & \zeta & \zeta^2 & 1 & \zeta^2 & \zeta^2 & \zeta & 1 \\ 1 & 1 & \zeta^2 & \zeta^2 & \zeta & 1 & \zeta & \zeta & \zeta^2 \\ 1 & \zeta & 1 & \zeta & \zeta & \zeta^2 & 1 & \zeta^2 & \zeta^2 \\ 1 & \zeta & \zeta^2 & 1 & \zeta^2 & \zeta^2 & \zeta & 1 & \zeta \\ 1 & \zeta^2 & \zeta^2 & \zeta & 1 & \zeta & \zeta & \zeta^2 & 1 \\ 1 & 1 & \zeta & \zeta & \zeta^2 & 1 & \zeta^2 & \zeta^2 & \zeta \\ 1 & \zeta^2 & 1 & \zeta^2 & \zeta^2 & \zeta & 1 & \zeta & \zeta \end{pmatrix}$$

# Hermitian codes over $\mathbb{F}_9$

$\mathcal{C}(9^H)$  is a group of order 192 with Molien series

$$\frac{\theta(t)}{(1-t^2)^2(1-t^4)^2(1-t^6)^3(1-t^8)(1-t^{12})}$$

where

$$\begin{aligned}\theta(t) := \quad & 1 + 3t^4 + 24t^6 + 74t^8 + 156t^{10} + 321t^{12} + 525t^{14} + 705t^{16} \\ & + 905t^{18} + 989t^{20} + 931t^{22} + 837t^{24} + 640t^{26} + 406t^{28} \\ & + 243t^{30} + 111t^{32} + 31t^{34} + 9t^{36} + t^{38},\end{aligned}$$

So the invariant ring of  $\mathcal{C}(9^H)$  has at least

$$\theta(1) + 9 = 6912 + 9 = 6921$$

generators and the maximal degree (=length of the code) is 38.  
What about Hamming weight enumerators ?

# Hermitian codes over $\mathbb{F}_9$

$$U := \text{ZU}(9^H) = \{x \in \mathbb{F}_9^* \mid x\bar{x} = x^4 = 1\} = (\mathbb{F}_9^*)^2$$

has 3 orbits on  $V = \mathbb{F}_9$ :

$$\{0\} = X_0, \{1, \alpha^2, \alpha^4, \alpha^6\} =: X_1, \{\alpha, \alpha^3, \alpha^5, \alpha^7\} =: X_2$$

$$\mathcal{C}^{(U)}(9^H) = \langle d_\varphi^{(U)} := \text{diag}(1, \zeta, \zeta^2), m_\alpha^{(U)} := \begin{pmatrix} 100 \\ 001 \\ 010 \end{pmatrix}, h^{(U)} := \frac{1}{3} \begin{pmatrix} 1 & 4 & 4 \\ 1 & 1 & -2 \\ -1 & 2 & 1 \end{pmatrix} \rangle$$

of order  $\frac{192}{4} = 48$  of which the invariant ring is a polynomial ring spanned by the  $U$ -symmetrized weight enumerators

$$q_2 = x_0^2 + 8x_1x_2, \quad q_4 = x_0^4 + 16(x_0x_1^3 + x_0x_2^3 + 3x_1^2x_2^2)$$

$$\begin{aligned} q_6 = & x_0^6 + 8(x_0^3x_1^3 + x_0^3x_2^3 + 2x_1^6 + 2x_2^6) \\ & + 72(x_0^2x_1^2x_2^2 + 2x_0x_1^4x_2 + 2x_0x_1x_2^4) + 320x_1^3x_2^3 \end{aligned}$$

of the three codes with generator matrices

$$\left[ \begin{array}{cc} 1 & \alpha \end{array} \right], \left[ \begin{array}{cccc} 1 & 1 & 1 & 0 \\ 0 & 1 & 2 & 1 \end{array} \right], \left[ \begin{array}{cccccc} 1 & 1 & 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 0 & 0 & 0 \\ 0 & \alpha & 2\alpha & 0 & 1 & 2 \end{array} \right].$$

# Hermitian codes over $\mathbb{F}_9$

Their Hamming weight enumerators are

$$r_2 = q_2(x, y, y) := x^2 + 8y^2,$$

$$r_4 = q_4(x, y, y) := x^4 + 32xy^3 + 48y^4,$$

$$r_6 = q_6(x, y, y) := x^6 + 16x^3y^3 + 72x^2y^4 + 288xy^5 + 352y^6.$$

The polynomials  $r_2, r_4$  and  $r_6$  generate the ring  $\text{Ham}(9^H)$  spanned by the Hamming weight enumerators of the codes of Type  $9^H$ .

$\text{Ham}(9^H) = \mathbb{C}[r_2, r_4] \oplus r_6\mathbb{C}[r_2, r_4]$  with the syzygy

$$r_6^2 = \frac{3}{4}r_2^4r_4 - \frac{3}{2}r_2^2r_4^2 - \frac{1}{4}r_4^3 - r_2^3r_6 + 3r_2r_4r_6.$$

Note that  $\text{Ham}(9^H)$  is **not** the invariant ring of a finite group.